Advanced Core in Algorithm Design #12
算法設計要論 第12回

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Schedule

Lec. #	Date	Topics
1	10/4	Introduction, Stable matching
2	10/11	Basics of Algorithm Analysis, Greedy Algorithms (1/2)
3	10/18	Greedy Algorithms (2/2)
4	10/25	Divide and Conquer (1/2)
5	11/1	Divide and Conquer (2/2)
6	11/8	Dynamic Programming (1/2)
7	11/15	Dynamic Programming (2/2)
—	11/22	Thursday Classes
8	11/29	Network Flow (1/2)
9	12/6	Network Flow (2/2)
10	12/13	NP and Computational Intractability
11	12/20	Approximation Algorithms (1/2)
12	12/27	Approximation Algorithms (2/2)
13	1/10	Randomized Algorithms

Approximation algorithm (repeated)

Definition

For a maximization problem "$\max f(x)$ s.t. $x \in X$", $0 \leq \alpha \leq 1$
a solution $x^* \in X$ is an α-approximation solution if $f(x^*) \geq \alpha \cdot \mathrm{OPT}$

Definition

For a minimization problem "$\min f(x)$ s.t. $x \in X$", $\alpha \geq 1$
a solution $x^* \in X$ is an α-approximation solution if $f(x^*) \leq \alpha \cdot \mathrm{OPT}$

Definition

An α-approximation algorithm is a polynomial-time algorithm that finds an α-approximation solution for any instance

Outline

Set Cover Problem

Problem

- Input: $U = \{e_1, \ldots, e_n\}$, $S_1, S_2, \ldots, S_m \subseteq U$
- Goal: minimize $|J|$ s.t. $J \subseteq [m]$ and $\bigcup_{j \in J} S_j = U$

This problem is **NP**-hard since the decision version is **NP**-complete

Example

- $U = \{e_1, e_2, e_3, e_4, e_5, e_6, e_7\}$
- $S_1 = \{e_1, e_2, e_3\}$
- $S_2 = \{e_1, e_4, e_5, e_7\}$
- $S_3 = \{e_3, e_6, e_7\}$
- $S_4 = \{e_4, e_5, e_6, e_7\}$

Greedy algorithm

Algorithm

1 $J \leftarrow \emptyset$;
2 **while** $\bigcup_{j \in J} S_j \subsetneq U$ **do**
3 Select $j^* \in [m]$ that maximizes $|S_{j^*} \setminus \bigcup_{j \in J} S_j|$;
4 $J \leftarrow J \cup \{j^*\}$;
5 **Return** J;

Example $J = \emptyset$

- $U = \{e_1, e_2, e_3, e_4, e_5, e_6, e_7\}$
- $S_1 = \{e_1, e_2, e_3\}$ → $S_1 \setminus \bigcup_{j \in J} S_j = \{e_1, e_2, e_3\}$
- $S_2 = \{e_1, e_4, e_5, e_7\}$ → $S_2 \setminus \bigcup_{j \in J} S_j = \{e_1, e_4, e_5, e_7\}$
- $S_3 = \{e_3, e_6, e_7\}$ → $S_3 \setminus \bigcup_{j \in J} S_j = \{e_3, e_6, e_7\}$
- $S_4 = \{e_4, e_5, e_6, e_7\}$ → $S_4 \setminus \bigcup_{j \in J} S_j = \{e_4, e_5, e_6, e_7\}$

Greedy algorithm

Algorithm

1 $J \leftarrow \emptyset$;
2 **while** $\bigcup_{j \in J} S_j \subsetneq U$ **do**
3 Select $j^* \in [m]$ that maximizes $|S_{j^*} \setminus \bigcup_{j \in J} S_j|$;
4 $J \leftarrow J \cup \{j^*\}$;
5 **Return** J;

Example $J = \{2\}$

- $U = \{e_1, e_2, e_3, e_4, e_5, e_6, e_7\}$
- $S_1 = \{e_1, e_2, e_3\}$ → $S_1 \setminus \bigcup_{j \in J} S_j = \{e_2, e_3\}$
- $S_2 = \{e_1, e_4, e_5, e_7\}$
- $S_3 = \{e_3, e_6, e_7\}$ → $S_3 \setminus \bigcup_{j \in J} S_j = \{e_3, e_6\}$
- $S_4 = \{e_4, e_5, e_6, e_7\}$ → $S_4 \setminus \bigcup_{j \in J} S_j = \{e_6\}$

Greedy algorithm

Algorithm

1 $J \leftarrow \emptyset$;
2 **while** $\bigcup_{j \in J} S_j \subsetneq U$ **do**
3 | Select $j^* \in [m]$ that maximizes $|S_{j^*} \setminus \bigcup_{j \in J} S_j|$;
4 | $J \leftarrow J \cup \{j^*\}$;
5 **Return** J;

Example $J = \{1, 2\}$

- $U = \{e_1, e_2, e_3, e_4, e_5, e_6, e_7\}$
- $S_1 = \{e_1, e_2, e_3\}$
- $S_2 = \{e_1, e_4, e_5, e_7\}$
- $S_3 = \{e_3, e_6, e_7\}$ → $S_3 \setminus \bigcup_{j \in J} S_j = \{e_6\}$
- $S_4 = \{e_4, e_5, e_6, e_7\}$ → $S_4 \setminus \bigcup_{j \in J} S_j = \{e_6\}$

Greedy algorithm

Algorithm

1 $J \leftarrow \emptyset$;
2 **while** $\bigcup_{j \in J} S_j \subsetneq U$ **do**
3 | Select $j^* \in [m]$ that maximizes $|S_{j^*} \setminus \bigcup_{j \in J} S_j|$;
4 | $J \leftarrow J \cup \{j^*\}$;
5 **Return** J;

Example $J = \{1, 2, 3\}$

- $U = \{e_1, e_2, e_3, e_4, e_5, e_6, e_7\}$
- $S_1 = \{e_1, e_2, e_3\}$
- $S_2 = \{e_1, e_4, e_5, e_7\}$
- $S_3 = \{e_3, e_6, e_7\}$
- $S_4 = \{e_4, e_5, e_6, e_7\} \longrightarrow S_4 \setminus \bigcup_{j \in J} S_j = \emptyset$

Analysis of the greedy algorithm

Theorem

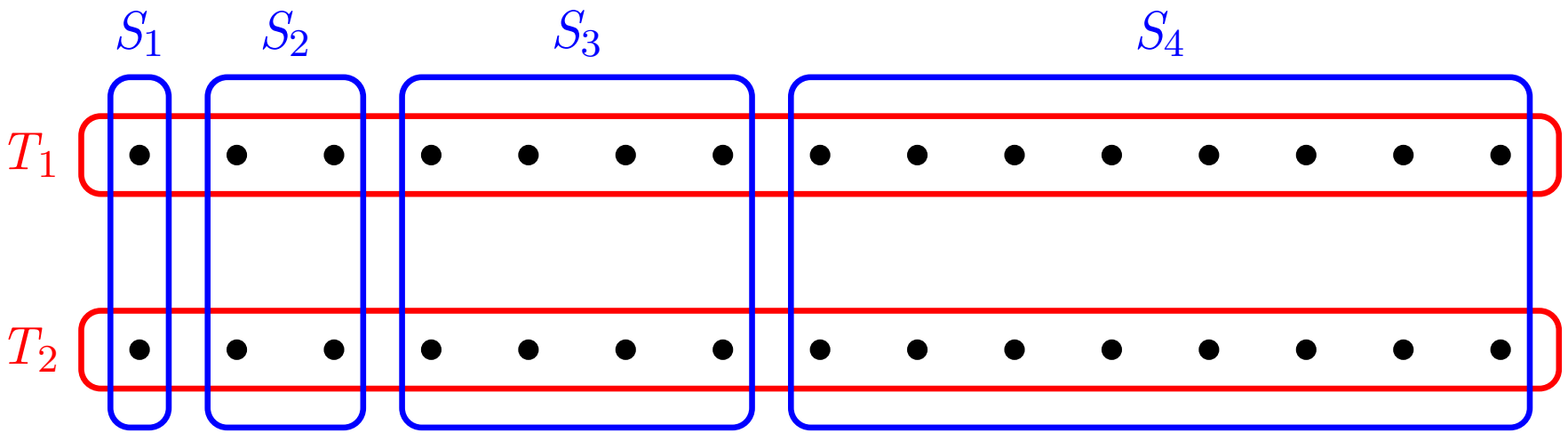
The greedy algorithm is an $\mathrm{O}(\log n)$-approximation algorithm

Proof

- $e_1, e_2, \ldots, e_n$: an order in which they were covered by the algorithm
- $p(e_k) :=$ (# of covered elements at the same time as e_k)
  $\geq \dfrac{n-k+1}{\mathrm{OPT}}$ ($\because$ at least $n-k+1$ elements are not covered before e_k is covered)
- $\mathrm{ALG} = \displaystyle\sum_{e_k \in U} \frac{1}{p(e_k)} \leq \sum_{e_k \in U} \frac{\mathrm{OPT}}{n-k+1} = \sum_{t=1}^{n} \frac{\mathrm{OPT}}{t} = H_n \cdot \mathrm{OPT} = O(\log n) \cdot \mathrm{OPT}$

Worst case example

The greedy algorithm outputs $\Omega(\log n)$-approx. solution for the following



Outline

Monotone Submodular Function

Definition

A function $f\colon 2^E \to \mathbb{R}$ is called

- monotone if $f(X) \leq f(Y)$ $(\forall X \subseteq \forall Y \subseteq E)$
- submodular if $f(X) + f(Y) \geq f(X \cup Y) + f(X \cap Y)$ $(\forall X, Y \subseteq E)$
- normalized if $f(\emptyset) = 0$

Example

- linear function: $f(X) = \sum_{i \in X} w_i$
- coverage function: $f(X) = \left|\bigcup_{i \in X} S_i\right|$
- matroid rank function: $f(X) = \max\{|X'| : X' \subseteq X,\ X' \in \mathcal{I}\}$
- joint entropy

Monotone Submodular Maximization

Problem

- Input: monotone submodular function $f\colon 2^E \to \mathbb{R}_+$ and $k \in \{1, \ldots, |E|\}$
- Goal: maximize $f(X)$ subject to $|X| \leq k$

Set cover problem $\leq_{\mathrm{P}}$ Monotone submodular maximization problem

→ Monotone submodular maximization problem is **NP**-hard

Example

- $U = \{1, 2, \ldots, 13\}$
- $S_1 = \{1, 2, 3\}$
- $S_2 = \{2, 3, 4, 5, 6\}$
- $S_3 = \{4, 5, 6, 7, 8, 9, 10\}$
- $S_4 = \{8, 9, 10, 11, 12, 13\}$
- $E = \{1, 2, 3, 4\}$
- $f(X) = \left|\bigcup_{i \in X} S_i\right|$
- $k = 2$

Greedy algorithm

1 Initially $S^{(0)} \leftarrow \emptyset$;
2 **for** $\ell \leftarrow 1, 2, \ldots, k$ **do**
3 Let $e^{(\ell)} \in \arg\max\{f(S^{(\ell-1)} \cup \{e\}) - f(S^{(\ell-1)}) \mid e \in E \setminus S^{(\ell-1)}\}$;
4 $S^{(\ell)} \leftarrow S^{(\ell-1)} \cup \{e^{(\ell)}\}$;
5 **Return** $S^{(k)}$;

Example

- $U = \{1, 2, \ldots, 13\}$
- $S_1 = \{1, 2, 3\}$
- $S_2 = \{2, 3, 4, 5, 6\}$
- $S_3 = \{4, 5, 6, 7, 8, 9, 10\}$
- $S_4 = \{8, 9, 10, 11, 12, 13\}$
- $E = \{1, 2, 3, 4\}$
- $f(X) = \left|\bigcup_{i \in X} S_i\right|$
- $k = 2$

Approximation ratio

Lemma 1

For all $T \subseteq T' \subseteq E$, $f(T') - f(T) \leq \sum_{e \in T' \setminus T} (f(T \cup \{e\}) - f(T))$

Lemma 2

$f(S^{(\ell)}) - f(S^{(\ell-1)}) \geq \frac{1}{|S^* \setminus S^{(\ell-1)}|}(f(S^*) - f(S^{(\ell-1)})) \geq \frac{1}{k}(f(S^*) - f(S^{(\ell-1)}))$

Theorem

$f(S^{(\ell)}) \geq \left(1 - \left(1 - \frac{1}{k}\right)^{\ell}\right) f(S^*)$

The greedy algorithm is a $(1 - 1/e)$-approximation algorithm because $f(S^{(k)}) \geq \left(1 - \left(1 - \frac{1}{k}\right)^{k}\right) f(S^*) \geq \left(1 - \frac{1}{e}\right) f(S^*)$ by the theorem

Proof of Lemma 1

Lemma 1

For all $T \subseteq T' \subseteq E$, $f(T') - f(T) \leq \sum_{e \in T' \setminus T} (f(T \cup \{e\}) - f(T))$

Proof

- Let $T' \setminus T = \{e_1, e_2, \ldots, e_m\}$
- $T_0 = T, T_1 = T \cup \{e_1\}, \ldots, T_\ell = T \cup \{e_1, \ldots, e_\ell\}, \ldots, T_m = T'$
- By submodularity, $f(T_{\ell-1} \cup \{e_\ell\}) - f(T_{\ell-1}) \leq f(T \cup \{e_\ell\}) - f(T)$
- By summing up the inequalities, we get the desired inequality

Proof of Lemma 2

Lemma 1

For all $T \subseteq T' \subseteq E$, $f(T') - f(T) \leq \sum_{e \in T' \setminus T} (f(T \cup \{e\}) - f(T))$

Lemma 2

$f(S^{(\ell)}) - f(S^{(\ell-1)}) \geq \frac{1}{|S^* \setminus S^{(\ell-1)}|}(f(S^*) - f(S^{(\ell-1)})) \geq \frac{1}{k}(f(S^*) - f(S^{(\ell-1)}))$

Proof

- For each $\ell \in \{1, \ldots, k\}$, we have

$$
\begin{aligned}
f(S^*) - f(S^{(\ell-1)}) &\leq f(S^* \cup S^{(\ell-1)}) - f(S^{(\ell-1)}) \\
&\overset{\text{by Lemma 1}}{\leq} \sum_{e \in S^* \setminus S^{(\ell-1)}} (f(S^{(\ell-1)} \cup \{e\}) - f(S^{(\ell-1)})) \\
&\leq |S^* \setminus S^{(\ell-1)}| \cdot \max_{e \in S^* \setminus S^{(\ell-1)}} (f(S^{(\ell-1)} \cup \{e\}) - f(S^{(\ell-1)})) \\
&\leq |S^* \setminus S^{(\ell-1)}| \cdot (f(S^{(\ell)}) - f(S^{(\ell-1)}))
\end{aligned}
$$

- The second inequality of Lemma 2 holds by $|S^* \setminus S^{(\ell-1)}| \leq k$

Proof of Theorem

Theorem

$$f(S^{(\ell)}) \geq \left(1 - \left(1 - \tfrac{1}{k}\right)^{\ell}\right) f(S^*)$$

Proof by induction

- Base step ($\ell = 0$): $f(S^{(0)}) = 0 = \left(1 - (1 - 1/k)^0\right) f(S^*)$
- Induction step ($\ell \geq 1$):

$$\begin{aligned}
f(S^{(\ell+1)}) &\geq f(S^{(\ell)}) + \frac{1}{k}\left(f(S^*) - f(S^{(\ell)})\right) \\
&= \left(1 - \frac{1}{k}\right) \cdot f(S^{(\ell)}) + \frac{1}{k} \cdot f(S^*) \\
&\geq \left(1 - \frac{1}{k}\right)\left(1 - \left(1 - \frac{1}{k}\right)^{\ell}\right) \cdot f(S^*) + \frac{1}{k} \cdot f(S^*) \\
&= \left(1 - \left(1 - \frac{1}{k}\right)^{\ell+1}\right) \cdot f(S^*)
\end{aligned}$$

by Lemma 2

induction hypothesis

Outline

Knpsack problem

Problem

- Input: items $E = \{1, 2, \ldots, n\}$ and a capacity $W \in \mathbb{Z}_+$
  item i has value $v_i \in \mathbb{Z}_+$ and size $w_i \in \mathbb{Z}_+$
- Goal: maximize $\sum_{i \in I} v_i$ subject to $I \subseteq \{1, 2, \ldots, n\}$ $\sum_{i \in I} w_i \leq W$

Recap: knapsack problem is **NP**-hard

Examples

$W = 16$

i	v_i	w_i
1	55	4
2	61	2
3	82	9
4	38	1
5	63	3

Knpsack problem

Problem

- Input: items $E = \{1, 2, \ldots, n\}$ and a capacity $W \in \mathbb{Z}_+$
  item i has value $v_i \in \mathbb{Z}_+$ and size $w_i \in \mathbb{Z}_+$
- Goal: maximize $\sum_{i \in I} v_i$ subject to $I \subseteq \{1, 2, \ldots, n\}$ $\sum_{i \in I} w_i \leq W$

Recap: knapsack problem is **NP**-hard

Examples

$W = 16$

i	v_i	w_i
1	55	4
2	61	2
3	82	9
4	38	1
5	63	3

optimal value is $61 + 82 + 38 + 63 = 244$

Continuous knapsack problem

(Integral) Knapsack problem

$$\max \quad \sum_{i \in E} v_i x_i \quad \text{s.t.} \quad \sum_{i \in E} w_i x_i \leq W, \quad x_i \in \{0, 1\} \ (\forall i \in E)$$

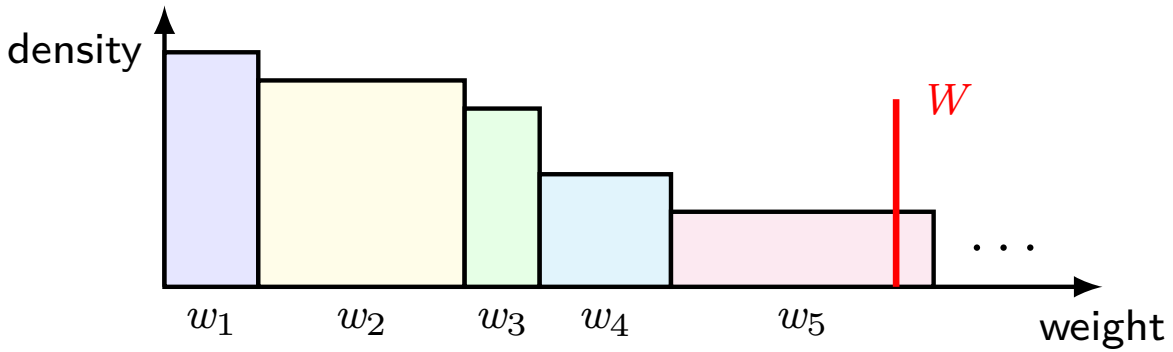
Continuous Knapsack problem

$$\max \quad \sum_{i \in E} v_i x_i \quad \text{s.t.} \quad \sum_{i \in E} w_i x_i \leq W, \quad x_i \in [0, 1] \ (\forall i \in E)$$

Observations

- $\mathrm{OPT}^{\mathsf{int}} \leq \mathrm{OPT}^{\mathsf{cont}}$
- fractional knapsack problem can be solved by a greedy algorithm

  descending order of their density (values per unit weight v_i/w_i)



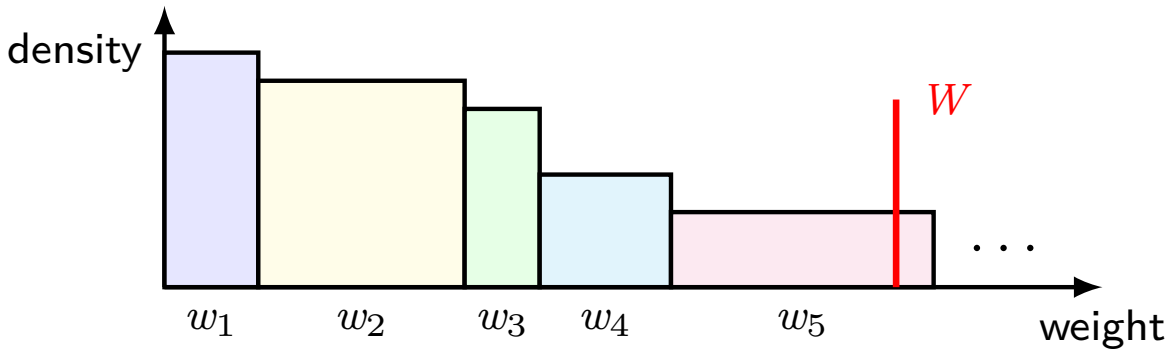
$1/2$-approximation algorithm

1. Sort items (with size $\leq W$) and relabel so that $\frac{v_1}{w_1} \geq \frac{v_2}{w_2} \geq \cdots \geq \frac{v_n}{w_n}$;
2. Let k be the index such that $\sum_{i=1}^{k-1} w_i \leq W < \sum_{i=1}^{k} w_i$;
3. Pick the better of $\{1, 2, \ldots, k-1\}$ and $\{k\}$;

Theorem

The above algorithm is $1/2$-approximation

- $\mathbf{x}^* = (\underbrace{1}_{1}, \ldots, \underbrace{1}_{k-1}, \underbrace{\frac{W - \sum_{i=1}^{k-1} w_i}{w_k}}_{k}, \underbrace{0}_{k+1}, \ldots, \underbrace{0}_{n})$ is optimal for cont. ver.
- $2 \max\left\{\sum_{i=1}^{k-1} v_i, v_k\right\} \geq \sum_{i=1}^{k} v_i \geq \sum_{i=1}^{n} v_i x_i^* = \mathrm{OPT}^{\mathsf{cont}} \geq \mathrm{OPT}^{\mathsf{int}}$

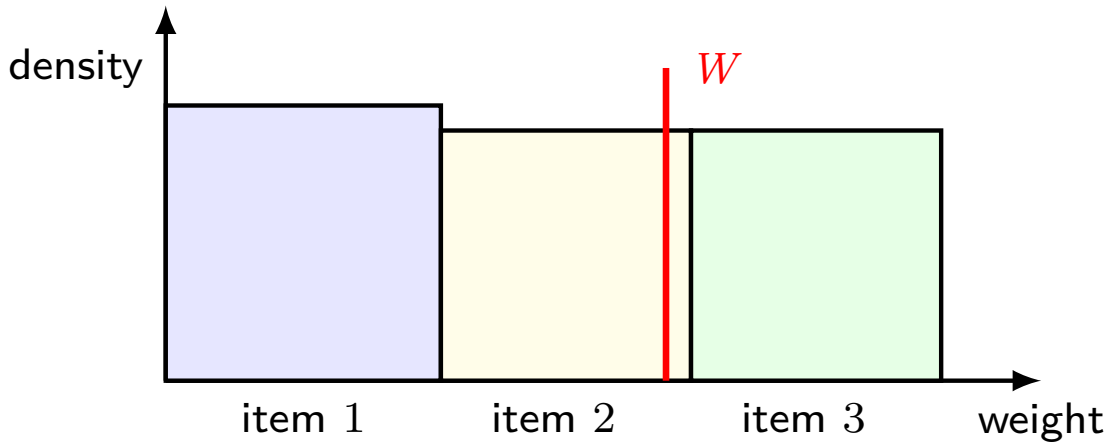


Worst case

Instance

$n = 3$ and $W = 2$

- item 1: $v_1 = (1+\epsilon)^2$, $w_1 = 1+\epsilon$
- item 2: $v_2 = 1$, $w_2 = 1$
- item 3: $v_3 = 1$, $w_3 = 1$



Analysis $\frac{\mathrm{ALG}}{\mathrm{OPT}} = \frac{(1+\epsilon)^2}{2} \to \frac{1}{2} \quad (\epsilon \to 0)$

PTAS and FPTAS (for maximization problems)

Definition: Polynomial-time approximation scheme (PTAS)

A PTAS is a $(1-\epsilon)$-approximation algorithm that runs in time polynomial in the problem size for any constant $\epsilon > 0$

Definition: Fully polynomial-time approximation scheme (FPTAS)

A FPTAS is a $(1-\epsilon)$-approximation algorithm that runs in time polynomial in both the problem size and $1/\epsilon$ for any $\epsilon > 0$

- A $(1-\epsilon)$-approximation algorithm that runs in $\mathrm{O}(n^{1/\epsilon})$ or $\mathrm{O}(n^{(1/\epsilon)^{1/\epsilon}})$ is PTAS but not FPTAS
- We will see PTAS and FPTAS for the knapsack problem

PTAS

Idea: guess the top-ℓ for value in the optimal solution

Algorithm

1 Sort items (with size $\leq W$) and relabel so that $\frac{v_1}{w_1} \geq \frac{v_2}{w_2} \geq \cdots \geq \frac{v_n}{w_n}$;
2 **foreach** *nonempty* $X \subseteq E$ *with* $|X| \leq \ell$ *and* $\sum_{i \in X} w_i \leq W$ **do**
3 $\quad S_X \leftarrow X$ and $v^* \leftarrow \min_{i \in X} v_i$;
4 $\quad$ **for** $i \leftarrow 1, 2, \ldots, n$ **do**
5 $\quad\quad$ **if** $i \notin X$, $v_i \leq v^*$, *and* $w(S_X) + w_i \leq W$ **then**
6 $\quad\quad\quad S_X \leftarrow S_X \cup \{i\}$;
7 **Return** the optimal solution among S_X;

Theorem

The above algorithm is $(1 - \frac{1}{\ell+1})$-approximation and runs in $\mathrm{O}(n^{\ell+1})$ time

- $\mathrm{O}(n^{\ell})$ possibilities of X imply $\mathrm{O}(n^{\ell+1})$ time $+\mathrm{O}(n \log n)$ time for sort
- By setting $\ell = \lceil 1/\epsilon \rceil - 1$, it is $(1 - \epsilon)$-approx. alg. that runs in $\mathrm{O}(n^{\lceil 1/\epsilon \rceil})$ time

Proof of approximation ratio

Theorem

The algorithm is $(1 - \frac{1}{\ell+1})$-approximation

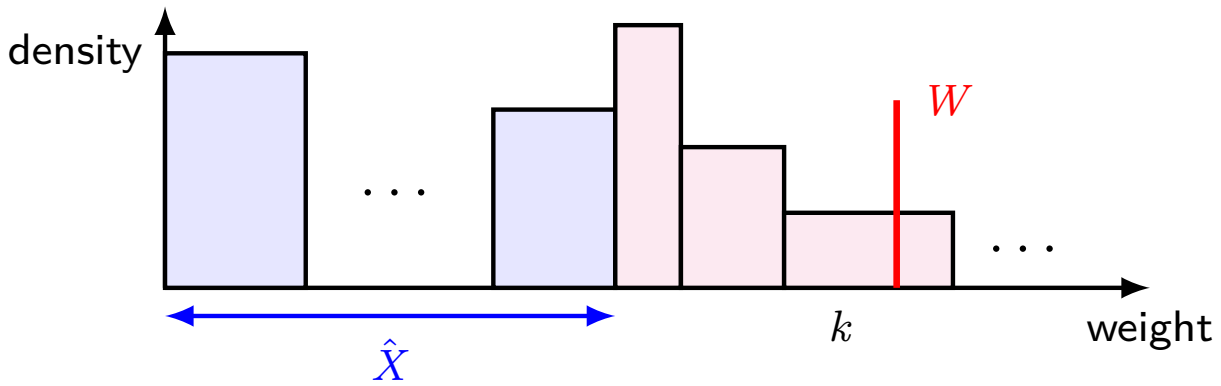
Proof

- Assumption: the optimal solution contains at least $\ell + 1$ items

  since otherwise the algorithm outputs the optimal solution

- Let X^* be the optimal solution and let $\hat{X}$ be the top-ℓ items in it
- $v(X^*) \leq v(S_{\hat{X}}) + \sum_{i=1}^{k} v_i \leq v(S_{\hat{X}}) + \sum_{i=1}^{k-1} v_i + \frac{v(X^*)}{\ell+1} \leq \mathrm{ALG} + \frac{v(X^*)}{\ell+1}$

  $\rightarrow \mathrm{ALG} \geq v(X^*) - \frac{v(X^*)}{\ell+1} = (1 - \frac{1}{\ell+1})\mathrm{OPT}$



Recap: Dynamic programming for knapsack problem

$$\mathrm{OPT}(k, w) = \max\{v(X) \mid X \subseteq \{1, 2, \ldots, k\},\ w(X) \leq w\}$$

Recursive formula

$$\mathrm{OPT}(k, w) = \begin{cases} 0 & \text{if } k = 0, \\ \mathrm{OPT}(k-1, w) & \text{if } w_k > w \\ \max\{\mathrm{OPT}(k-1, w),\ \mathrm{OPT}(k-1, b-w_k) + v_k\} & \text{otherwise} \end{cases}$$

Compute for $k = 0, 1, \ldots, n$ and $w = 0, 1, \ldots, W$ → $\mathrm{O}(nW)$ time

Example $(w_i, v_i) = (4, 55), (2, 61), (9, 82), (1, 38), (3, 63),\ W = 16$

$k \backslash w$	0	1	2	3	4	5	6	7	8	9	10	11	12	13	14	15	16
0	0	0	0	0	0	0	0	0	0	0	0	0	0	0	0	0	0
1	0	0	0	0	55	55	55	55	55	55	55	55	55	55	55	55	55
2	0	0	61	61	61	61	116	116	116	116	116	116	116	116	116	116	116
3	0	0	61	61	61	61	116	116	116	116	116	143	143	143	143	198	198
4	0	38	61	99	99	99	116	154	154	154	154	154	181	181	181	198	236
5	0	38	61	99	101	124	162	162	162	179	217	217	217	217	217	244	244

Another dynamic programming for knapsack problem

$$\mathrm{OPT}(k, v) = \min\{w(X) \mid X \subseteq \{1, 2, \ldots, k\},\ v(X) = v\}$$

Recursive formula

$$\mathrm{OPT}(k, v) = \begin{cases} 0 & \text{if } v = 0, \\ +\infty & \text{else if } k = 0, \\ \min\{\mathrm{OPT}(k-1, v),\ \mathrm{OPT}(k-1, v-v_k) + w_k\} & \text{else if } v \geq v_k, \\ \mathrm{OPT}(k-1, v) & \text{otherwise} \end{cases}$$

Compute for $k = 0, 1, \ldots, n$ and $v = 0, 1, \ldots, \sum_{i=1}^{n} v_i$ → $\mathrm{O}(n^2 \max_i v_i)$ time

Example $(w_i, v_i) = (55, 4), (61, 2), (82, 6), (38, 1), (63, 3),\ W = 160$

$k \backslash v$	0	1	2	3	4	5	6	7	8	9	10	11	12	13	14	15	16
0	0	∞	∞	∞	∞	∞	∞	∞	∞	∞	∞	∞	∞	∞	∞	∞	∞
1	0	∞	∞	∞	55	∞	∞	∞	∞	∞	∞	∞	∞	∞	∞	∞	∞
2	0	∞	61	∞	55	∞	116	∞	∞	∞	∞	∞	∞	∞	∞	∞	∞
3	0	∞	61	∞	55	∞	82	∞	143	∞	137	∞	198	∞	∞	∞	∞
4	0	38	61	99	55	93	82	120	143	181	137	175	198	236	∞	∞	∞
5	0	38	61	63	55	93	82	118	143	145	137	175	198	200	238	261	299

FPTAS

Idea: scale down the value of every item into $0, 1, \ldots, \ell$

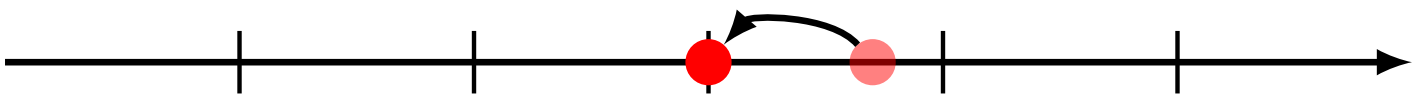
Algorithm

1. Let $v_{\max} = \max_{i'=1}^{n} v_{i'}$;
2. For each item i, let $\hat{v}_i = \left\lfloor \frac{v_i}{v_{\max}} \cdot \ell \right\rfloor$ (delete i with $w_i > W$);
3. Compute the optimal solution for $\left((w_1, \hat{v}_1), \ldots, (w_n, \hat{v}_n); W\right)$ by the dynamic programming and output it;

Theorem

The above algorithm is $(1 - \frac{n}{\ell})$-approximation and runs in $\mathrm{O}(n^2\ell)$ time

- The running time is $\mathrm{O}(n^2\ell)$ since $\max_i \hat{v}_i = \ell$
- By setting $\ell = \left\lfloor \frac{n}{\epsilon} \right\rfloor$, it is $(1 - \epsilon)$-approx. alg. that runs in $\mathrm{O}(\frac{n^3}{\epsilon})$ time



Proof of approximation ratio

Theorem

The algorithm is $(1 - \frac{n}{\ell})$-approximation

Proof

- Let X^* be the optimal solution and let X be the output of the algorithm
- $\frac{v_i}{v_{\max}} \cdot \ell - 1 < \hat{v}_i = \left\lfloor \frac{v_i}{v_{\max}} \cdot \ell \right\rfloor \leq \frac{v_i}{v_{\max}} \cdot \ell$
- Thus, we have

$$\sum_{i \in X} v_i \geq \sum_{i \in X} \hat{v}_i \cdot \frac{v_{\max}}{\ell} \geq \frac{v_{\max}}{\ell} \sum_{i \in X^*} \hat{v}_i \geq \frac{v_{\max}}{\ell} \sum_{i \in X^*} \left(\frac{v_i}{v_{\max}} \ell - 1 \right)$$

$$= \sum_{i \in X^*} v_i - \frac{v_{\max}}{\ell} |X^*| \geq \left(1 - \frac{|X^*|}{\ell}\right) \sum_{i \in X^*} v_i \geq \left(1 - \frac{n}{\ell}\right) \sum_{i \in X^*} v_i$$

X is optimal for $\hat{v}_i$

$v_{\max} \leq \sum_{i \in X^*} v_i$